# An Efficient Algorithm for Skeptical Preferred Acceptance in Dynamic Argumentation Frameworks

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### **ABSTRACT ARGUMENTATION**

An (*abstract*) *argumentation framework* (*AF*) is a pair  $\langle A, \Sigma \rangle$ , where *A* is a set of *arguments* and  $\Sigma \subseteq A \times A$  is a set of *attacks*.

- It allows representing dialogues, making decisions, and handling inconsistency;
- An AF can be viewed as a direct graph, whose nodes are arguments and whose edges are attacks.

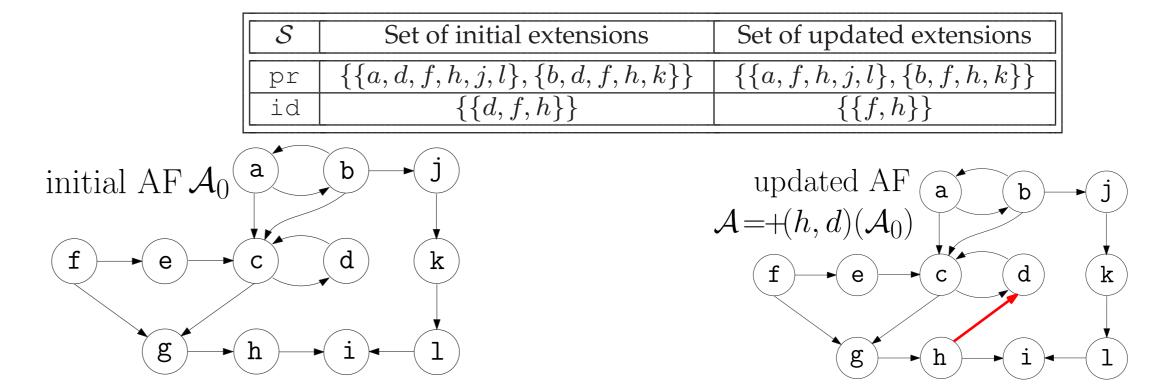
## **SEMANTICS FOR AFS**

An argumentation semantics specifies the criteria for identifying "reasonable" sets of arguments, called extensions.

- A *preferred extension* (pr) is a maximal (w.r.t.  $\subseteq$ ) admissible set.
- An *ideal extension* (id) is the biggest (w.r.t.

# **DYNAMIC ARGUMENTATION FRAMEWORKS**

- An argumentation framework models a temporary situation as **new arguments and attacks** can be added/removed to take into account new available knowledge.
- The set of arguments skeptically accepted under the preferred semantics may change if we update an initial AF  $A_0$  by adding/removing arguments/attacks. For instance, the skeptical acceptance under the preferred semantics of goal argument d is *true* in  $A_0$  but false in the updated AF  $\mathcal{A} = +(h, d)(\mathcal{A}_0)$  obtained from  $\mathcal{A}_0$  by adding attack (h, d). This is due to the change of the set of the preferred extensions.



- Should we recompute the skeptical acceptance of updated AFs from scratch?

 $\subseteq$ ) admissible set which is contained in every preferred extension.

An argument is **skeptically accepted** under the preferred semantics iff it belongs to every preferred extension.

#### **UPDATES**

An *update* u for an AF  $A_0$  consists in modifying  $A_0$  into an AF A by adding or removing arguments or attacks.

- + (a, b) (resp. -(a, b)) denotes the addition (resp. deletion) of an attack (a, b);
- $u(\mathcal{A}_0)$  means applying  $u = \pm(a, b)$  to  $\mathcal{A}_0$ ;
- multiple (attacks) updates can be simulated by a single attack update.

#### CONTRIBUTIONS

We show that computing a small portion of the input AF is sufficient to determine the skeptical acceptance of a goal argument in the updated AF.

We introduce SPA, an incremental algorithm for computing the Skeptical Preferred Acceptance of a goal within a dynamic AF. It consists of the following main steps:

- Identify a sub-AF called *context-based* AF on the basis of updates and additional information provided by the ideal extension.
- Give as input the context-based AF to an external (non-incremental) solver to compute (i) the skeptical preferred acceptance of the goal argument, and (ii) the ideal extension for the updated AF.

We provide a thorough experimental analysis showing the effectiveness of our approach.

10<sup>7</sup>

10<sup>6</sup>

10<sup>5</sup>

10<sup>4</sup>

10<sup>3</sup>

10<sup>2</sup>

10<sup>1</sup>

10<sup>0</sup>

 $10^{4}$ 

 $10^{3}$ 

10

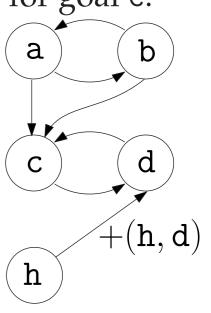
10

 $10^{0}$ 

 $10^{-1}$ 

 $10^{-1}$ 

Context-based AF for goal c:

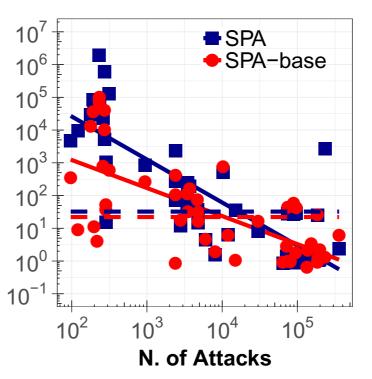


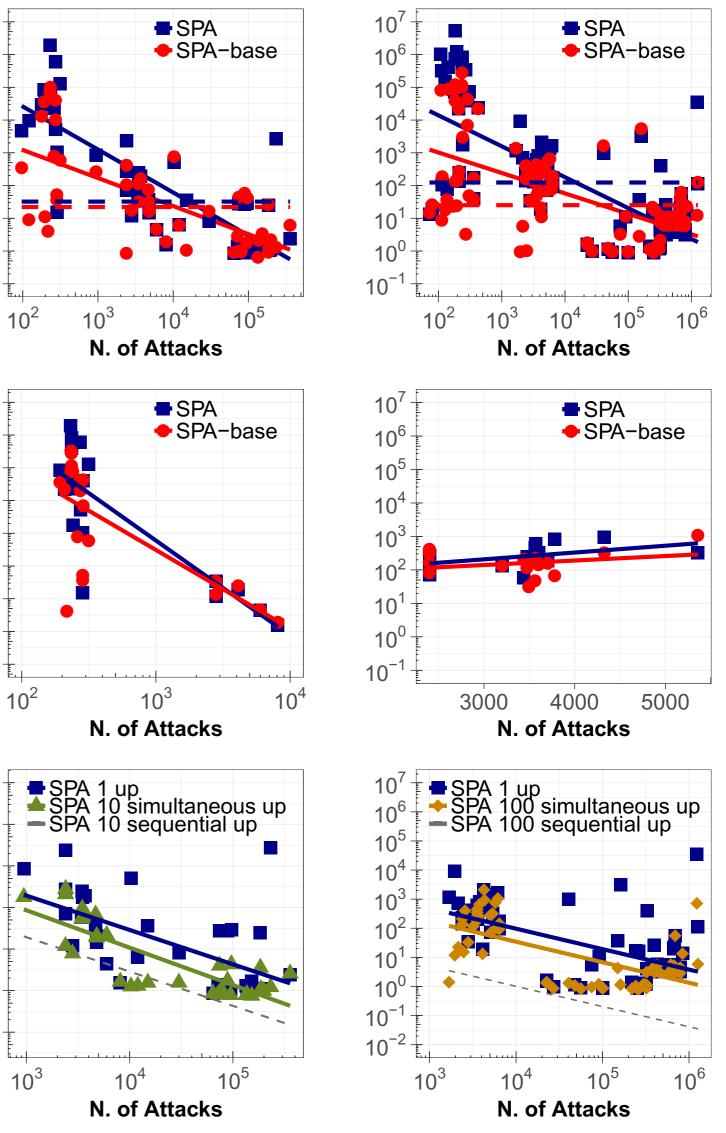
#### **EXPERIMENTS**

#### Datasets: ICCMA'17 benchmarks.

For each AF in the dataset, we compared the performance of our technique with that of *ArgSemSAT*, the solver that won the last ICCMA competition for the computational task DS-pr: Given an AF, determine the skeptical preferred acceptance of a given argument.

**Results:** The figures report the improvements (running time of the competitor over running time of our approach) of SPA and SPA-base versus the number of attacks. SPA-base is a version of SPA not using the ideal extension.





- Considering the averages of the improvements, SPA and SPA-base turn out to be 5 and 4 orders of magnitude faster than *ArgSemSAT*, respectively. However, as this can be skewed by extremely large values of improvements (e.g.  $10^6$ ), we also considered the medians of improvements for SPA (32 on A2, 134 on A3) and SPA-base (27 on A2, 40 on A3) (see dashed line), which confirm the significance of the gain in efficiency.
- SPA is generally faster than SPA-base, except for a few AFs whose initial ideal extension is empty.
- The performance gets worse when the ratio between the size of the contextbased AF and that of the initial AF becomes very large because of the increasing density of the initial AFs.
- For sets of updates, results show that SPA remains faster than the competitor even when 10 or 100 updates are performed simultaneously.
- Finally, applying updates simultaneously is faster than applying them sequentially.